1971 Turing Award Lecture

Generality in Artificial Intelligence

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The Turing Award Lecture given in 1971 by John McCarthy was never published. The postscript that follows, written by the author in 1986, endeavors to reflect the flavor of the original, as well as to comment in the light of development over the past 15 years.

Postscript

My 1971 Turing Award Lecture was entitled "Generality in Artificial Intelligence." The topic turned out to have been overambitious in that I discovered that I was unable to put my thoughts on the subject in a satisfactory written form at that time. It would have been better to have reviewed previous work rather than attempt something new, but such wasn't my custom at that time.

I am grateful to the ACM for the opportunity to try again. Unfortunately for our science, although perhaps fortunately for this project, the problem of generality in artificial intelligence (AI) is almost as unsolved as ever, although we now have many ideas not available in 1971. This paper relies heavily on such ideas, but it is far from a full 1986 survey of approaches for achieving generality. Ideas are discussed

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at a length proportional to my familiarity with them rather than according to some objective criterion.

It was obvious in 1971 and even in 1958 that AI programs suffered from a lack of generality. It is still obvious, and now there are many more details. The first gross symptom is that a small addition to the idea of a program often involves a complete rewrite beginning with the data structures. Some progress has been made in modularizing data structures, but small modifications of the search strategies are even less likely to be accomplished without rewriting.

Another symptom is that no one knows how to make a general database of common sense knowledge that could be used by any program that needed the knowledge. Along with other information, such a database would contain what a robot would need to know about the effects of moving objects around, what a person can be expected to know about his family, and the facts about buying and selling. This doesn't depend on whether the knowledge is to be expressed in a logical language or in some other formalism. When we take the logic approach to AI, lack of generality shows up in that the axioms we devise to express common sense knowledge are too restricted in their applicability for a general common sense database. In my opinion, getting a language for expressing general common sense knowledge for inclusion in a general database is the key problem of generality in AI.

Here are some ideas for achieving generality proposed both before and after 1971. I repeat my disclaimer of comprehensiveness.

Representing Behavior by Program

Friedberg [7, 8] discussed a completely general way of representing behavior and provided a way of learning to improve it. Namely, the behavior is represented by a computer program and learning is accomplished by making random modifications to the program and testing the modified program. The Friedberg approach was successful in learning only how to move a single bit from one memory cell to another, and its scheme of rewarding instructions involved in successful runs by reducing their probability of modification was shown by Simon [24] to be inferior to testing each program thoroughly and completely scrapping any program that wasn't perfect. No one seems to have attempted to follow up the idea of learning by modifying whole programs.

The defect of the Friedberg approach is that while representing behaviors by programs is entirely general, modifying behaviors by small modifications to the programs is very special. A small conceptual modification to a behavior is usually not represented by a small modification to the program, especially if machine language programs are used and any one small modification to the text of a program is considered as likely as any other.

It might be worth trying something more analogous to genetic evolution; duplicates of subroutines would be made, some copies would

be modified and others left unchanged. The learning system would then experiment whether it was advantageous to change certain calls of the original subroutine to calls of the modified subroutine. Most likely even this wouldn't work unless the relevant small modifications of behavior were obtainable by calls to slightly modified subroutines. It would probably be necessary to provide for modifications to the number of arguments of subroutines.

While Friedberg's problem was learning from experience, all schemes for representing knowledge by program suffer from similar difficulties when the object is to combine disparate knowledge or to make programs that modify knowledge.

The General Problem Solver (GPS) and Its Successor

One kind of generality in AI comprises methods for finding solutions that are independent of the problem domain. Allen Newell, Herbert Simon, and their colleagues and students pioneered this approach and continue to pursue it.

Newell et al. first proposed GPS in 1957 [18]. The initial idea was to represent problems of some general class as problems of transforming one expression into another by means of a set of allowed rules. It was even suggested in [20] that improving GPS could be thought of as a problem of this kind. In my opinion, GPS was unsuccessful as a general problem solver because problems don't take this form in general and because most of the knowledge needed for problem solving and achieving goals is not simply representable in the form of rules for transforming expressions. However, GPS was the first system to separate the problem-solving structure of goals and subgoals from the particular domain.

If GPS had worked out to be really general, perhaps the Newell and Simon predictions about rapid success for AI would have been realized. Newell's current candidate [22] for general problem representation is SOAR, which, as I understand it, is concerned with transforming one state to another, where the states need not be represented by expressions.

Production Systems

The first production systems were done by Newell and Simon in the 1950s, and the idea was written up in [21]. A kind of generality is achieved by using the same goal-seeking mechanism for all kinds of problems, changing only the particular productions. The early production systems have grown into the current proliferation of expert system shells.

Production systems represent knowledge in the form of facts and rules, and there is almost always a sharp syntactic distinction between the rules. The facts usually correspond to ground instances of logical

formulas, that is, they correspond to predicate symbols applied to constant expressions. Unlike logic-based systems, these facts contain no variables or quantifiers. New facts are produced by inference, observation, and user input. Variables are reserved for rules, which usually take a pattern-action form. Rules are put in the system by the programmer or "knowledge engineer" and in most systems cannot arise via the action of the system. In exchange for accepting these limitations, the production system programmer gets a relatively fast program.

Production system programs rarely use fundamental knowledge of the domain. For example, MYCIN [2] has many rules about how to infer which bacterium is causing an illness based on symptoms and the result of laboratory tests. However, its formalism has no way of expressing the fact that bacteria are organisms that grow within the body. In fact, MYCIN has no way of representing processes occurring in time, although other production systems can represent processes at about the level of the situation calculus to be described in the next section.

The result of a production system pattern match is a substitution of constants for variables in the pattern part of the rule. Consequently, production systems do not infer general propositions. For example, consider the definition that a container is sterile if it is sealed against entry by bacteria, and all the bacteria in it are dead. A production system (or a logic program) can only use this fact by substituting particular bacteria for the variables. Thus it cannot reason that heating a sealed container will sterilize it given that a heated bacterium dies, because it cannot reason about the unenumerated set of bacteria in the container. These matters are discussed further in [14].

Representing Knowledge in Logic

It seemed to me in 1958 that small modifications in behavior are most often representable as small modifications in beliefs about the world, and this requires a system that represents beliefs explicitly.

> If one wants a machine to be able to discover an abstraction, it seems most likely that the machine must be able to represent this abstraction in some relatively simple way. [11, p. 78]

The 1960 idea for increasing generality was to use logic to express facts in a way independent of the way the facts might subsequently be used. It seemed then and still seems that humans communicate mainly in declarative sentences rather than in programming languages for good objective reasons that will apply whether the communicator is a human, a creature from Alpha Centauri, or a computer program. Moreover, the advantages of declarative information also apply to internal representation. The advantage of declarative information is one of generality. The fact that when two objects collide they make a noise may be used in particular situations to make a noise, to avoid making noise, to explain a noise, or to explain the absence of noise. (I guess those cars didn't collide, because while I heard the squeal of brakes, I didn't hear a crash.)

Once one has decided to build an AI system that represents information declaratively, one still has to decide what kind of declarative language to allow. The simplest systems allow only constant predicates applied to constant symbols, for example, *on(Block1, Block2)*. Next, one can allow arbitrary constant terms, built from function symbols constants and predicate symbols, for example, *location(Block1) = top(Block2)*. Prolog databases allow arbitrary Horn clauses that include free variables, for example, $P(x, y) \wedge Q(y, z) \supset R(x, z)$, expressing the Prolog in standard logical notation. Beyond that lies full first-order logic, including both existential and universal quantifiers and arbitrary first-order formulas. Within first-order logic, the expressive power of a theory depends on what domains the variables are allowed to range. Important expressive power comes from using set theory, which contains expressions for sets of any objects in the theory.

Every increase in expressive power carries a price in the required complexity of the reasoning and problem-solving programs. To put it another way, accepting limitations on the expressiveness of one's declarative information allows simplification of the search procedures. Prolog represents a local optimum in this continuum, because Horn clauses are medium expressive but can be interpreted directly by a logical problem solver.

One major limitation that is usually accepted is to limit the derivation of new facts to formulas without variables, that is, to substitute constants for variables and then do propositional reasoning. It appears that most human daily activity involves only such reasoning. In principle, Prolog goes slightly beyond this, because the expressions found as values of variables by Prolog programs can themselves involve free variables. However, this facility is rarely used except for intermediate results.

What can't be done without more of predicate calculus than Prolog allows is universal generalization. Consider the rationale of canning. We say that a container is sterile if it is sealed and all the bacteria in it are dead. This can be expressed as a fragment of a Prolog program as follows:

```
sterile(X) :- sealed(X), not alive-bacterium(Y, X).
alive-bacterium(Y, X) :- in(Y, X), bacterium(Y), alive(Y).
```

However, a Prolog program incorporating this fragment directly can sterilize a container only by killing each bacterium individually and would require that some other part of the program successively generate the names of the bacteria. It cannot be used to discover or rationalize canning—sealing the container and then heating it to kill all the bacteria at once. The reasoning rationalizing canning involves the use of quantifiers in an essential way.

My own opinion is that reasoning and problem-solving programs will eventually have to allow the full use of quantifiers and sets and have strong enough control methods to use them without combinatorial explosion.

While the 1958 idea was well received, very few attempts were made to embody it in program in the immediately following years, the main one being Black's Harvard Ph.D. dissertation of 1964. I spent most of my time on what I regarded as preliminary projects, mainly LISP. My main reason for not attempting an implementation was that I wanted to learn how to express common sense knowledge in logic first. This is still my goal. I might be discouraged from continuing to pursue it if people pursuing nonlogical approaches were having significant success in achieving generality.

McCarthy and Hayes [12] made the distinction between epistemological and heuristic aspects of the AI problem and asserted that generality is more easily studied epistemologically. The distinction is that the epistemology is completed when the facts available have as a consequence that a certain strategy is appropriate to achieve the goal, whereas the heuristic problem involves the search that finds the appropriate strategy.

Implicit in [11] was the idea of a general-purpose, common sense database. The common sense information possessed by humans would be written as logical sentences and included in the database. Any goal-seeking program could consult the database for the facts needed to decide how to achieve its goal. Especially prominent in the database would be facts about the effects of actions. The much studied example is the set of facts about the effects of a robot trying to move objects from one location to another. This led in the 1960s to the *situation calculus* [12], which was intended to provide a way of expressing the consequences of actions independent of the problem.

The basic formalism of the situation calculus is

$$s' = result(e, s),$$

which asserts that s' is the situation that results when event e occurs in situation s. Here are some situation calculus axioms for moving and painting blocks.

Qualified Result-of-Action Axioms

$$\forall x\ l\ s.clear(top(x), s) \wedge clear(l, s) \wedge \neg\ tooheavy(x) \supset loc(x, result(move(x, l), s)) = l.$$

$$\forall x\ c\ s.color(x, result(paint(x, c), s)) = c.$$

Frame Axioms

$$\forall x\ y\ l\ s.color(y, result(move(x, l), s)) = color(y, s).$$

$$\forall x\ y\ l\ s.y \neq x \supset loc(y, result(move(x, l), s)) = loc(y, s).$$

$$\forall x\ y\ c\ s.loc(x, result(paint(y, c), s)) = loc(x, s).$$

$$\forall x\ y\ c\ s.y \neq x \supset color(x, result(paint(y, c), s)) = color(x, s).$$

Notice that all qualifications to the performance of the actions are explicit in the premises and that statements (called frame axioms) about

what doesn't change when an action is performed are explicitly included. Without those statements it wouldn't be possible to infer much about *result*(*e2*, *result*(*e1*, *s*)), since we wouldn't know whether the premises for the event *e2* to have its expected result were fulfilled in *result*(*e1*, *s*).

Notice further that the situation calculus applies only when it is reasonable to reason about discrete events, each of which results in a new total situation. Continuous events and concurrent events are not covered.

Unfortunately, it wasn't very feasible to use the situation calculus in the manner proposed, even for problems meeting its restrictions. In the first place, using general-purpose theorem provers made the programs run too slowly, since the theorem provers of 1969 [9] had no way of controlling the search. This led to STRIPS [6], which reduced the use of logic to reasoning within a situation. Unfortunately, the STRIPS formalizations were much more special than full situation calculus. The facts that were included in the axioms had to be delicately chosen in order to avoid the introduction of contradictions arising from the failure to delete a sentence that wouldn't be true in the situation that resulted from an action.

Nonmonotonicity

The second problem with the situation calculus axioms is that they were again not general enough. This was the *qualification problem,* and a possible way around it wasn't discovered until the late 1970s. Consider putting an axiom in a *common sense database* asserting that birds can fly. Clearly the axiom must be *qualified* in some way since penguins, dead birds, and birds whose feet are encased in concrete can't fly. A careful construction of the axiom might succeed in including the exceptions of penguins and dead birds, but clearly we can think up as many additional exceptions, like birds with their feet encased in concrete, as we like. Formalized nonmonotonic reasoning (see [4], [15]–[17], and [23]) provides a formal way of saying that a bird can fly unless there is an abnormal circumstance and of reasoning that only the abnormal circumstances whose existence follows from the facts being taken into account will be considered.

Nonmonotonicity has considerably increased the possibility of expressing general knowledge about the effects of events in the situation calculus. It has also provided a way of solving the *frame problem,* which constituted another obstacle to generality that was already noted in [12]. The frame problem (the term has been variously used, but I had it first) occurs when there are several actions available, each of which changes certain features of the situation. Somehow it is necessary to say that an action changes only the features of the situation to which it directly refers. When there is a fixed set of actions and features, it can be explicitly stated which features are unchanged by an action, even though it may take a lot of axioms. However, if we imagine that

additional features of situations and additional actions may be added to the database, we face the problem that the axiomatization of an action is never completed. McCarthy [16] indicates how to handle this using *circumscription,* but Lifschitz [10] has shown that circumscription needs to be improved and has made proposals for this.

Here are some situation calculus axioms for moving and painting blocks using circumscription taken from [16].

Axioms about Locations and the Effects of Moving Objects

$$\forall x\ e\ s.\neg ab(aspect1(x, e, s)) \supset loc(x, result(e, s)) = loc(x, s).$$
$$\forall x\ l\ s.ab(aspect1(x, move(x, l), s)).$$
$$\forall x\ l\ s.\neg ab(aspect3(x, l, s)) \supset loc(x, result(move(x, l), s)) = l.$$

Axioms about Colors and Painting

$$\forall x\ e\ s.\neg ab(aspect2(x, e, s)) \supset color(x, result(e, s)) = color(x, s).$$
$$\forall x\ c\ s.ab(aspect2(x, paint(x, c), s)).$$
$$\forall x\ c\ s.\neg ab(aspect4(x, c, s)) \supset color(x, result(paint(x, c), s)) = c.$$

This treats the qualification problem, because any number of conditions that may be imagined as preventing moving or painting can be added later and asserted to imply the corresponding *ab aspect* . . . . It treats the frame problem in that we don't have to say that moving doesn't affect colors and painting locations.

Even with formalized nonmonotonic reasoning, the general common sense database still seems elusive. The problem is writing axioms that satisfy our notions of incorporating the general facts about a phenomenon. Whenever we tentatively decide on some axioms, we are able to think of situations in which they don't apply and a generalization is called for. Moveover, the difficulties that are thought of are often ad hoc like that of the bird with its feet encased in concrete.

Reification

Reasoning about knowledge, belief, or goals requires extensions of the domain of objects reasoned about. For example, a program that does backward chaining on goals uses them directly as sentences: *on*(*Block1*, *Block2*); that is, the symbol *on* is used as a predicate constant of the language. However, a program that wants to say directly that *on*(*Block1*, *Block2*) should be postponed until *on*(*Block2*, *Block3*) has been achieved needs a sentence like *precedes*(*on*(*Block2*, *Block3*), *on*(*Block1*, *Block2*)), and if this is to be a sentence of first-order logic, then the symbol *on* must be taken as a function symbol, and *on*(*Block1*, *Block2*) regarded as an object in the first-order language.

This process of making objects out of sentences and other entities is called *reification.* It is necessary for expressive power but again leads to complications in reasoning. It is discussed in [13].

Formalizing the Notion of Context

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Whenever we write an axiom, a critic can say that the axiom is true only in a certain context. With a little ingenuity the critic can usually devise a more general context in which the precise form of the axiom doesn't hold. Looking at human reasoning as reflected in language emphasizes this point. Consider axiomatizing "on" so as to draw appropriate consequences from the information expressed in the sentence, "The book is on the table." The critic may propose to haggle about the precise meaning of "on," inventing difficulties about what can be between the book and the table or about how much gravity there has to be in a spacecraft in order to use the word "on" and whether centrifugal force counts. Thus we encounter Socratic puzzles over what the concepts mean in complete generality and encounter examples that never arise in life. There simply isn't a most general context.

Conversely, if we axiomatize at a fairly high level of generality, the axioms are often longer than is convenient in special situations. Thus humans find it useful to say, "The book is on the table," omitting reference to time and precise identifications of what book and what table. This problem of how general to be arises whether the general common sense knowledge is expressed in logic, in program, or in some other formalism. (Some people propose that the knowledge is internally expressed in the form of examples only, but strong mechanisms using analogy and similarity permit their more general use. I wish them good fortune in formulating precise proposals about what these mechanisms are.)

A possible way out involves formalizing the notion of context and combining it with the circumscription method of nonmonotonic reasoning. We add a context parameter to the functions and predicates in our axioms. Each axiom makes its assertion about a certain context. Further axioms tell us that facts are inherited by more restricted context unless exceptions are asserted. Each assertion is also nonmonotonically assumed to apply in any particular more general context, but there again are exceptions. For example, the rules about birds flying implicitly assume that there is an atmosphere to fly in. In a more general context this might not be assumed. It remains to determine how inheritance to more general contexts differs from inheritance to more specific contexts.

Suppose that whenever a sentence p is present in the memory of a computer, we consider it as in a particular context and as an abbreviation for the sentence $holds(p, C)$, where C is the name of a context. Some contexts are very specific, so that Watson is a doctor in the context of Sherlock Holmes stories and a baritone psychologist in a tragic opera about the history of psychology.

There is a relation $c1 \leq c2$ meaning that context $c2$ is more general than context $c1$. We allow sentences like $holds(c1 \leq c2, c0)$ so that

even statements relating contexts can have contexts. The theory would not provide for any "most general context" any more than Zermelo - Frankel set theory provides for a most general set.

A logical system using contexts might provide operations of *entering* and *leaving* a context yielding what we might call *ultranatural deduction* allowing a sequence of reasoning like

$$
\begin{array}{l}
holds(p, C) \\
ENTER\ C \\
p \\
\quad \vdots \\
q \\
LEAVE\ C \\
holds(q, C).
\end{array}
$$

This resembles the usual logical natural deduction systems, but for reasons beyond the scope of this lecture, it is probably not correct to regard contexts as equivalent to sets of assumptions — not even infinite sets of assumptions.

All this is unpleasantly vague, but it's a lot more than could be said in 1971.

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Categories and Subject Descriptors:

I.2.3 [**Artificial Intelligence**]: Deduction and Theorem Proving — *logic programming;* I.2.4 [**Artificial Intelligence**]: Knowledge Representation Formalisms and Methods — *representation languages;* I.2.6 [**Artificial Intelligence**]: Learning — *concept learning*

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